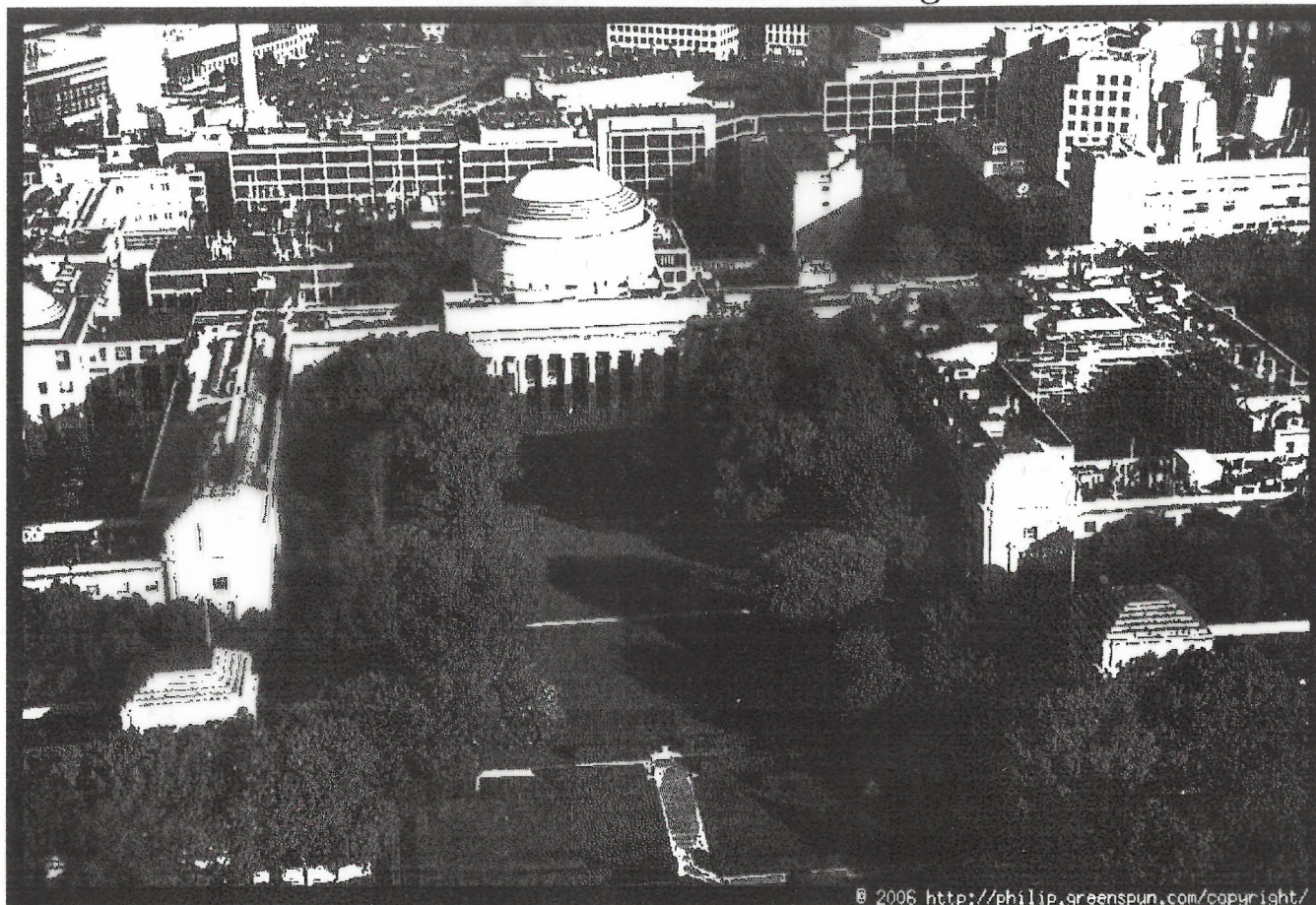


Build a Digital Clock from the 1980's

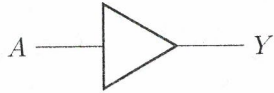
Lecture 2A: Combinatorial Logic



© 2006 <http://philip.greenspun.com/copyright/>

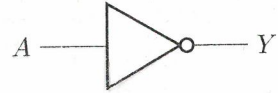
Original photo courtesy of Philip Greenspun (<http://philip.greenspun.com/>)
Edited using the GNU Image Manipulation Program (<https://www.gimp.org/>)

YES Gate



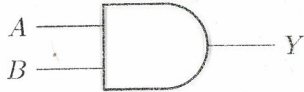
A	Y
0	0
1	1

NOT Gate



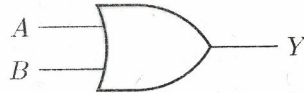
A	Y
0	1
1	0

AND Gate



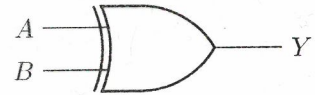
A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

OR Gate



A	B	Y
0	0	0
0	1	1
1	0	1
1	1	1

XOR Gate



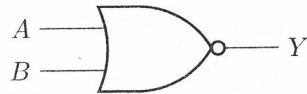
A	B	Y
0	0	0
0	1	1
1	0	1
1	1	0

NAND Gate



A	B	Y
0	0	1
0	1	1
1	0	1
1	1	0

NOR Gate



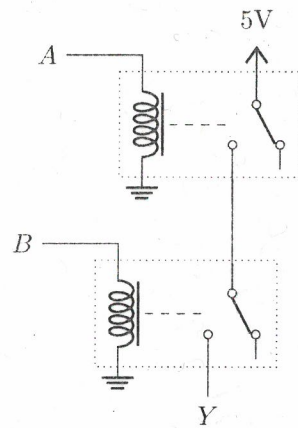
A	B	Y
0	0	1
0	1	0
1	0	0
1	1	0

XNOR Gate

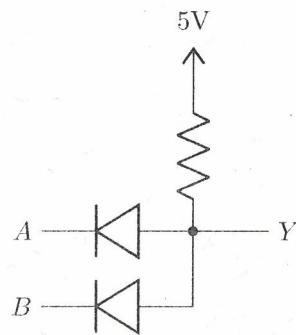


A	B	Y
0	0	1
0	1	0
1	0	0
1	1	1

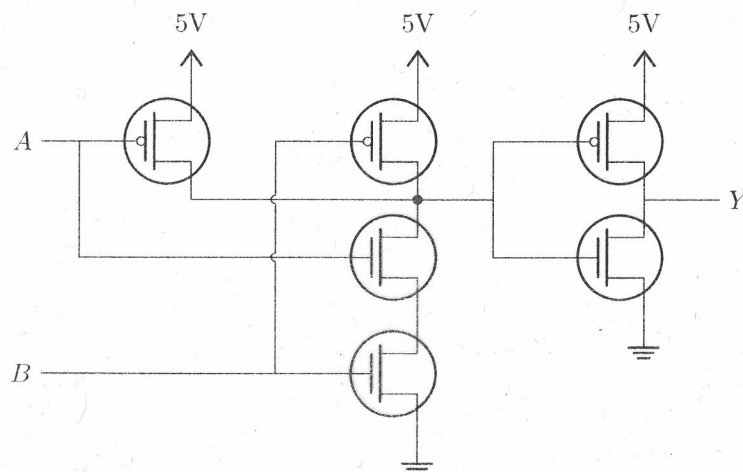
Relay AND Gate:



DRL AND Gate:



CMOS AND Gate:



Numerical Representations

Decimal:

6873

$$(6 \times 10^3) + (8 \times 10^2) + (7 \times 10^1) + (3 \times 10^0)$$

Sexagesimal:



$$\begin{aligned} &(1 \times 60^2) + (54 \times 60^1) + (33 \times 60^0) \\ &(1 \times 3600) + (54 \times 60) + (33 \times 1) \\ &3600 + 3240 + 33 = 6873 \end{aligned}$$

Binary:

1101011011001

$$\begin{aligned} &(1 \times 2^{12}) + (1 \times 2^{11}) + (1 \times 2^9) + (1 \times 2^7) + (1 \times 2^6) + (1 \times 2^4) + (1 \times 2^3) + (1 \times 2^0) \\ &4096 + 2048 + 512 + 128 + 64 + 16 + 8 + 1 \end{aligned}$$

$$\begin{array}{r} 172 \\ + 253 \\ \hline 425 \end{array}$$

$$\begin{array}{r}
 \begin{array}{c}
 \text{1} \quad \text{1} \quad \text{1} \quad \text{1} \quad \text{1} \\
 \text{0} \text{ } \text{1} \text{ } \text{1} \text{ } \text{1} \text{ } \text{0} \\
 + \quad \text{0} \text{ } \text{1} \text{ } \text{0} \text{ } \text{1} \text{ } \text{1} \\
 \hline
 \text{1} \text{ } \text{1} \text{ } \text{0} \text{ } \text{0} \text{ } \text{1} \\
 \text{16} \text{ } \text{8} \text{ } \text{4} \text{ } \text{2} \text{ } \text{1}
 \end{array}
 \end{array}$$

The diagram shows a binary addition of two 6-bit numbers. The first number is 01110 (decimal 14) and the second is 01011 (decimal 11). The sum is 11001 (decimal 25). The weights for each bit position are 16, 8, 4, 2, and 1.

MSB

	b	
a	0	1
0	0	0
1	0	1

ab

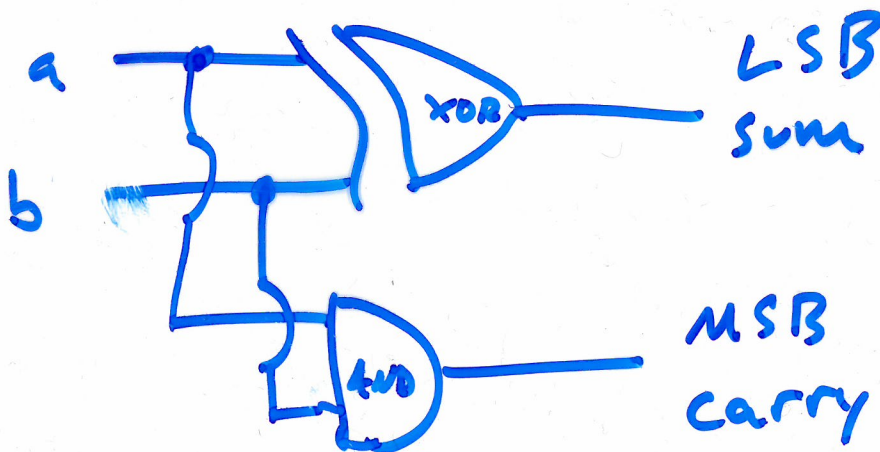
LSB

	b	
a	0	1
0	0	1
1	1	0

$a \oplus b$

$a+b$

"half-adder"



MSB
"carry"

c \ a, b	00	01	11	10
0	0	1	1	0
1	0	1	1	1

$$bc + ab + ac$$

3 AND, 2 OR

a	b	c	sum
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

$$ab + c(b + a)$$

$$ab + c((a \oplus b) + ab)$$

$$\underline{a}b + c(a \oplus b) + \cancel{ab}$$

$$ab + c(a \oplus b)$$

LSB

"sum"

c \ a, b	00	01	11	10
0	0	1	0	1
1	1	0	1	0

$$c\bar{a}\bar{b} + b\bar{a}\bar{c} + a\bar{b}\bar{c} + abc$$

8 ANDs + 3 OR's + 3 NOTs

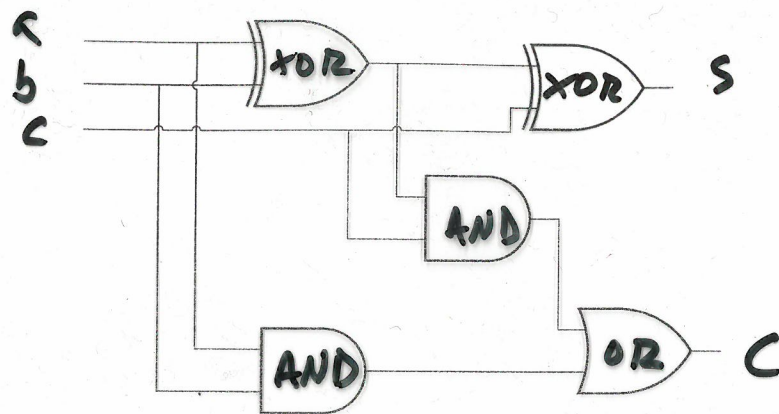
14

$$c(\underline{ab + \bar{a}\bar{b}}) + \bar{c}(\bar{a}b + \bar{b}a)$$

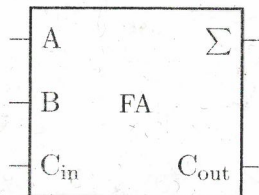
$$\underline{a \oplus b}$$

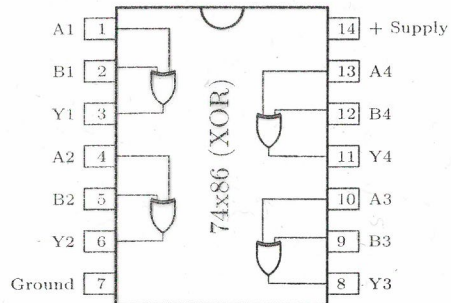
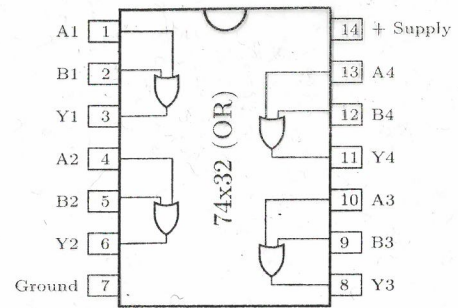
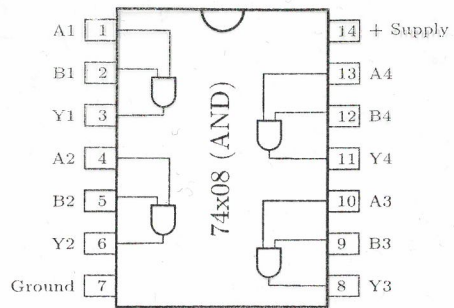
$$c(\overline{a \oplus b}) \quad \bar{c}(a \oplus b)$$

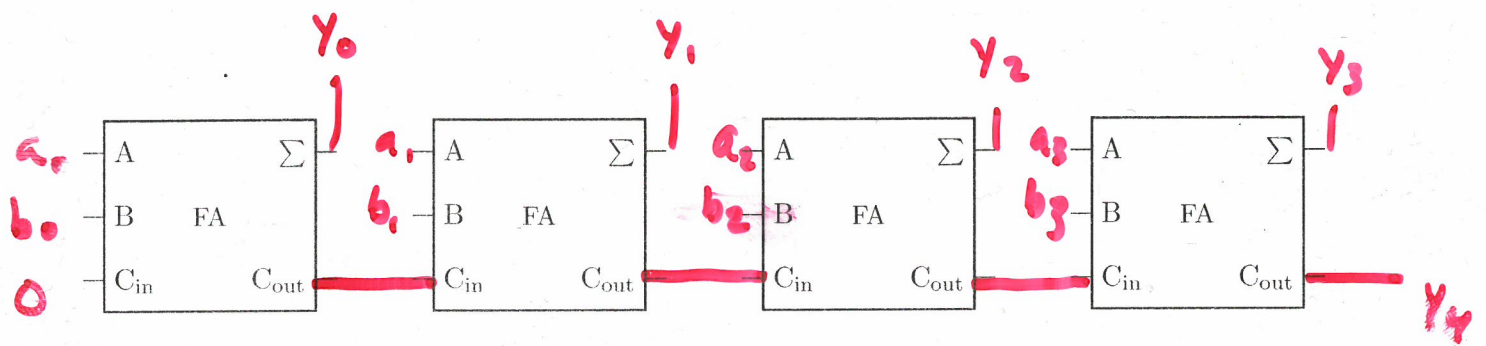
$$(a \oplus b) \oplus c$$



"full adder"







add

$a+b$

Numerical Representations Revisited

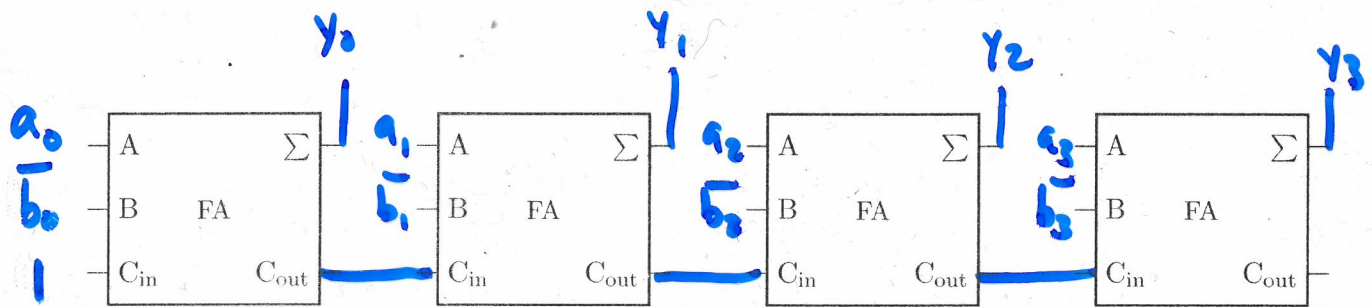
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	-8
1001	9	-7
1010	10	-6
1011	11	-5
1100	12	-4
1101	13	-3
1110	14	-2
1111	15	-1

4 → -4
0100 → 1100

$$2^N - x$$

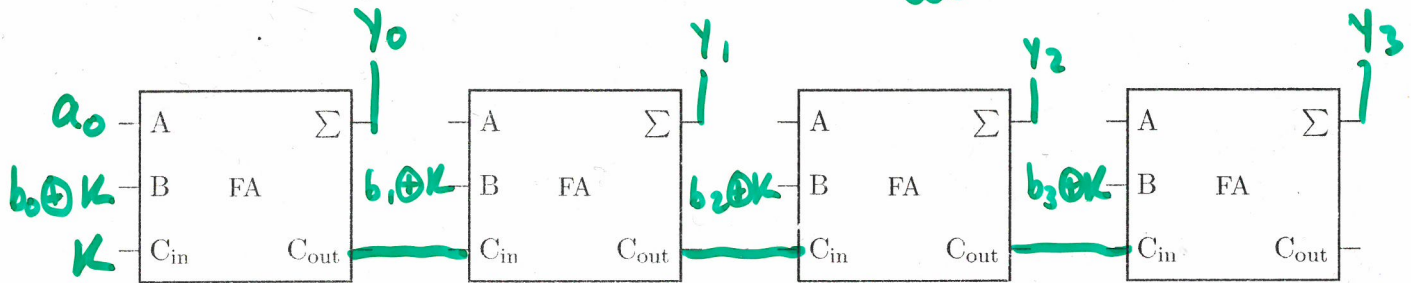
$$2^N - 1 - x + 1$$

$$\sim x + 1$$



Subtract $a-b$

K is a separate input: subtract if 1,
add if 0



$$y = a + b \text{ if } K = 0 \text{ else } a - b$$