



Here is a truth table.
Start with Sum of Products (S.O.P.)
and simplify it

Part I

a	b	c
0	0	0
0	1	1
1	0	1
1	1	1

S.O.P.:

$$c = \bar{a}b + ab + a\bar{b}$$

$$\text{hmm } c = \bar{a}b + a(\bar{b} + b)$$

$$= \bar{a}b + a(1)$$

$$= \bar{a}b + a$$

... hmm what if instead..

$$c = \bar{a}b + ab + a\bar{b}$$

$$(\bar{a} + a)b + a\bar{b}$$

$$c = (1)b + a\bar{b}$$

$$= b + a\bar{b}$$

... hmm WTF?

WE MUST USE IDEMPOTENCE

$a = a + a$ or $a = a \cdot a$
so retry.

$$c = \bar{a}b + ab + a\bar{b}$$

$$= \bar{a}b + \bar{a}b + ab + a\bar{b}$$

$$b(\bar{a} + a) + a(b + \bar{b})$$

$$c = b + a$$

noice

yes this is just
OR... but let's note
sure Boolean Algebra
can help us reduce
to simplified result.
should reduce to $c = a + b$

COULD HAVE ALSO
GENERATED A
TRUTH TABLE

a	b	c	y
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

Keep this in mind since
it does quickly indicate to
us this is a



(1960s EE
boss. could
smoke in office
then)

Build a circuit with
output y and inputs
a, b, c.

y should be high
if a is on or if
b is on with c off
or if c is on
or if b is off with
a on.

Just written out logic

$$y = a + b\bar{c} + c + \bar{b}a$$

$$\text{then } y = a(1 + \bar{b}) + b\bar{c} + c$$

$$y = a + b\bar{c} + c$$

Build
circuit



$$y = a + b\bar{c} + c = b + c$$

$$= a + b\bar{c} + cb + c\bar{b} + c$$

$$= a + b(\bar{c} + c) + c(1 + \bar{b})$$

$$= a + b(1) + c(1)$$

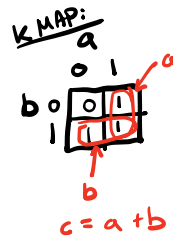
$$y = a + b + c$$

IS THERE A
BETTER / MORE RELIABLE
WAY TO REDUCE??
YES. KARNAUGH
MAPS!!

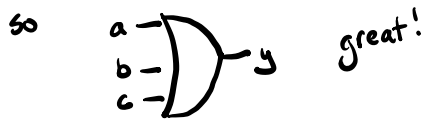
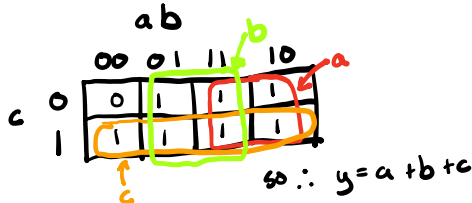
KARNAUGH MAP
ASIDE...
HIGHER DIMENSIONAL
TRUTH TABLE

OR

a	b	c
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



As applied to our $y = a + b\bar{c} + c + \bar{b}a$

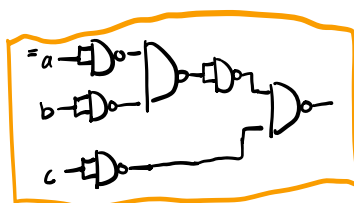
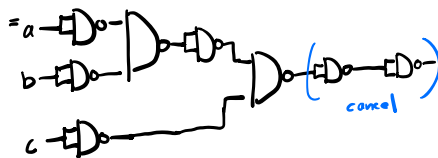
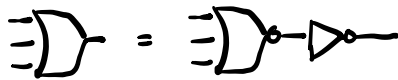


BUT WHAT IF?
....



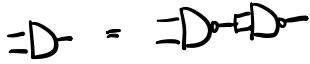
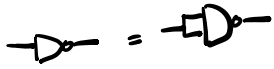
We only have 7400
Quad NAND chips.
You have to build
with that.

NEED deMorgan

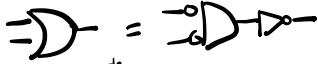


stuff like this
happened a lot.
or sometimes you'd have
left over gates so you'd
try to use those gates
to avoid bringing in more
chips to the design

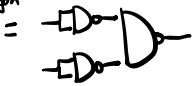
All logic can be represented
as sum of products. (S.O.P.)
S.O.P. has + • -
make those from $\Rightarrow$ you're good



$$A+B = \overline{\overline{A+B}} = \overline{\overline{A} \cdot \overline{B}} \rightarrow \overline{\overline{A} \cdot \overline{B}}$$

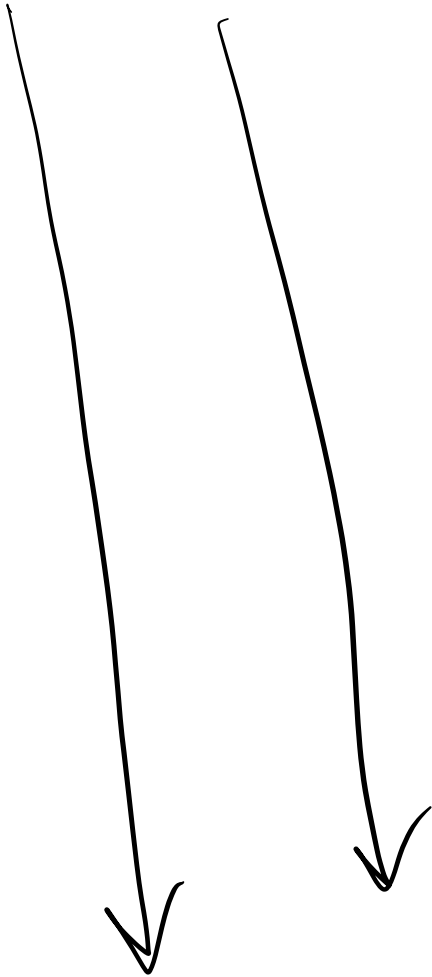


de
morgan



Yay! could do some
for NOR. I'll leave
that up to you.

Let's build some more legit logic...



A = 2 bits B = 2 bits

$$Y = A > B$$

$$A_1 \bar{A}_0 \bar{B}_1 (\bar{B}_0 + B_0)$$

$$Y = \bar{A}_1 A_0 \bar{B}_1 \bar{B}_0 + (A_1 \bar{A}_0 \bar{B}_1 \bar{B}_0 + A_1 \bar{A}_0 B_1 B_0) + (A_1 A_0 \bar{B}_1 \bar{B}_0 + A_1 A_0 \bar{B}_1 B_0) + A_1 A_0 B_1 \bar{B}_0$$

$$A_1 A_0 (\bar{B}_1 \bar{B}_0 + \bar{B}_1 B_0) \quad A_1 A_0 \bar{B}_0 (\bar{B}_1 + B_1)$$

$$A_1 A_0 (\bar{B}_1 (\bar{B}_0 + B_0))$$

$$A_1 A_0 \bar{B}_1$$

$$Y = \bar{A}_1 A_0 \bar{B}_1 \bar{B}_0 + A_1 \bar{A}_0 \bar{B}_1$$

Idempotent

$$+ A_1 A_0 \bar{B}_1 + A_1 A_0 B_1 \bar{B}_0$$

$$A_1 \bar{B}_1 \quad A_0 \bar{B}_0 (\bar{A}_1 B_1 + A_1 B_1)$$

$$A_0 \bar{B}_1 \bar{B}_0 (\bar{A}_1)$$

$$A_1 A_0 B_1 \bar{B}_0 + A_1 \bar{B}_1$$

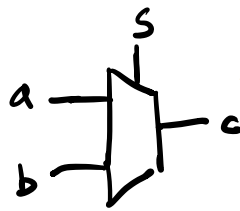
A ₁	A ₀	B ₁	B ₀	Y
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	1
1	0	0	1	1
1	0	1	0	0
1	0	1	1	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0

KARNAUGH MAP

		A ₁ A ₀		
		00	01	11
B ₁ B ₀	00	0	1	1
	01	0	0	1
	11	0	0	0
	10	0	0	1

NOT SURE!!!

Now MAKE
MULTIPLEXER



if (s)
c = a
else
c = b

Like a
solid state SPDT
RELAY

STEP 1: TRUTH TABLE

a	b	s	c
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

KMAP:

	ab			
	00	01	11	10
s	0	0	1	1
1	0	0	1	1

just two
for full coverage

$$c = as + b\bar{s}$$

$$c = \bar{a}b\bar{s} + a\bar{b}s + ab\bar{s} + abs$$

sum of products

could simplify from
here if we wanted.

